I.-C. Lin, Assistant Professor. Textbook: Operating System Principles 7ed

## CHAPTER 2: SYSTEM STRUCTURES

# Chapter 2: System Structures

- Operating System Services
- User Operating System Interface
- System Calls
- Types of System Calls
- System Programs
- Operating System Design and Implementation
- Operating System Structure
- Virtual Machines
- Operating System Generation
- System Boot

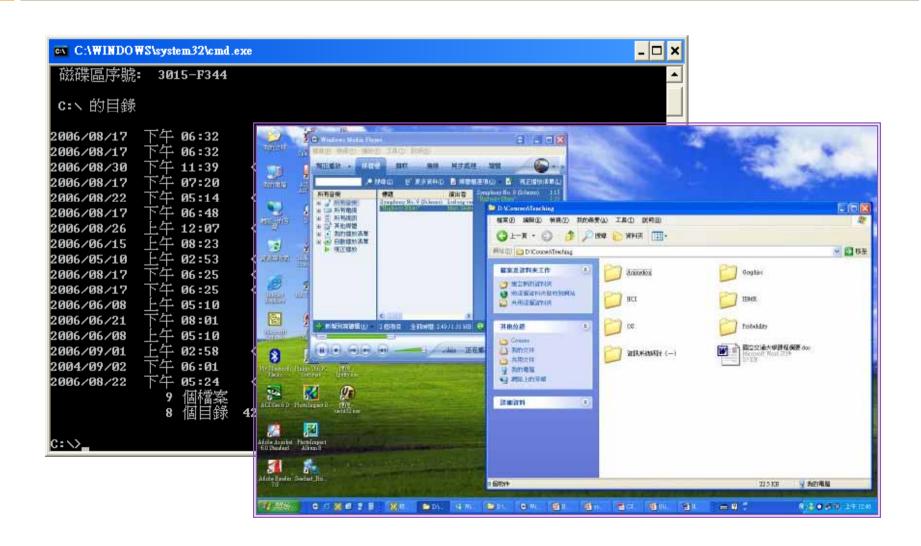
# Objectives

- To describe the services an operating system provides to users, processes, and other systems
- To discuss the various ways of structuring an operating system
- To explain how operating systems are installed and customized and how they boot

# **Operating System Services**

- Services for the convenience of the programmers (or users)
  - User interface
    - Varies between Command-Line (CLI), Graphics User Interface (GUI), Batch
  - Program execution
    - Load a program into memory and to run that program, end execution, either normally or abnormally (indicating error)
  - I/O operations
    - A running program may require I/O, which may involve a file or an I/O device.

## CLI vs. GUI



# Operating System Services (cont.)

### File-system manipulation

Read and write, create and delete, search files and directories, list file Information, permission management.

#### Communications

- Processes may exchange information, on the same computer or over a network
- May be via shared memory or through message passing

#### Error detection

- Error may occur in the CPU and memory hardware, in I/O devices, in user program
- Should take the appropriate action to ensure correct and consistent computing
- Debugging facilities for users to efficiently use the system

# Operating System Services (Cont.)

## For ensuring efficient system operations

#### Resource allocation

- allocating resources to multiple users or multiple jobs running at the same time
- Many types of resources.

## Accounting

- To keep track of which users use how much and what kinds of computer resources
- Statistics for improving computing services

### Protection and security

- Protection involves ensuring that all access to system resources is controlled
- Security of the system from outsiders requires user authentication, extends to defending external I/O devices from invalid access attempts

# User Operating System Interface - CLI

## Command-line interpreter allows direct command entry

- Sometimes implemented in kernel, sometimes by systems program
- Sometimes multiple flavors implemented shells
- Primarily fetches a command from user and executes it
  - Sometimes commands built-in, sometimes just names of programs
    - If the latter, adding new features doesn't require shell modification

## User Operating System Interface - GUI

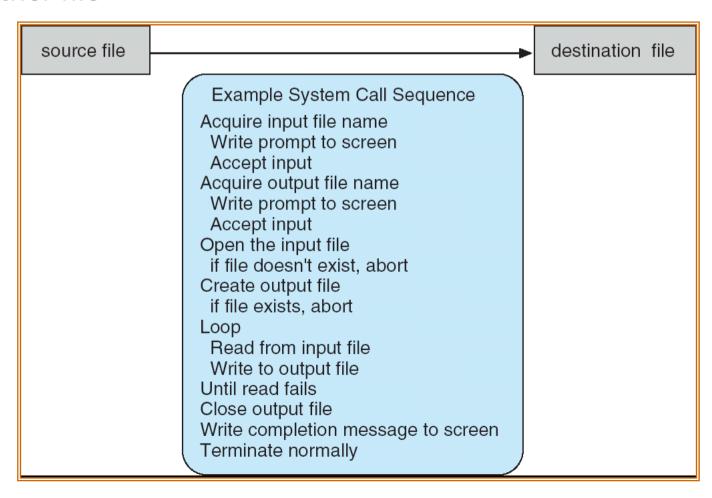
- User-friendly desktop metaphor interface
  - Usually mouse, keyboard, and monitor
  - Icons represent files, programs, actions, etc.
  - Various mouse buttons over objects in the interface cause various actions
- Many systems now include both CLI and GUI interfaces
  - Microsoft Windows is GUI with CLI "command" shell
  - Apple Mac OS X as "Aqua" GUI interface with UNIX kernel underneath and shells available
  - Solaris is CLI with optional GUI interfaces

# System Calls

- Programming interface to the services provided by the OS
- Typically written in a high-level language (C or C++)
- Mostly accessed by programs via a high-level Application
   Program Interface (API) rather than direct system call use
  - Three most common APIs are Win32 API for Windows, POSIX API for POSIX-based systems (including virtually all versions of UNIX, Linux, and Mac OS X), and Java API for the Java virtual machine (JVM)
- Why do we use APIs rather than system calls?

# Example of System Calls

 System call sequence to copy the contents of one file to another file



# **Example of Standard API**

- Consider the ReadFile() function in the Win32 API—a function for reading from a file
- A description of the parameters passed to ReadFile()
  - HANDLE file—the file to be read
  - LPVOID buffer—a buffer where the data will be read into and written from
  - DWORD bytesToRead—the number of bytes to be read into the buffer
  - LPDWORD bytesRead—the number of bytes read during the last read
  - LPOVERLAPPED ovl—indicates if overlapped I/O is being used

```
return value

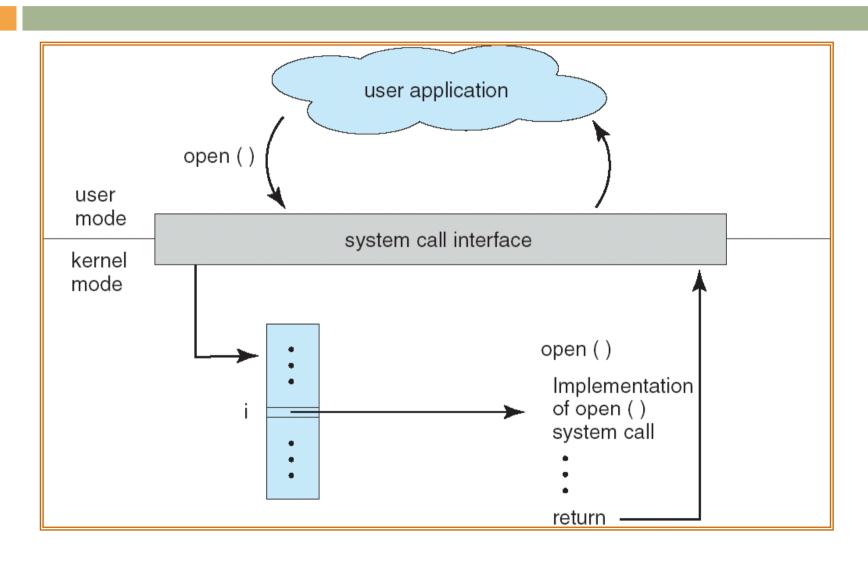
BOOL ReadFile c (HANDLE file,
LPVOID buffer,
DWORD bytes To Read,
LPDWORD bytes Read,
function name

LPOVERLAPPED ovl);
```

# System Call Implementation

- The system call interface
  - invokes intended system call in OS kernel and returns status of the system call and any return values
- The caller need know nothing about how the system call is implemented
  - Just needs to obey API and understand what OS will do as a result call
  - Most details of OS interface hidden from programmer by API
    - Managed by run-time support library (set of functions built into libraries included with compiler)

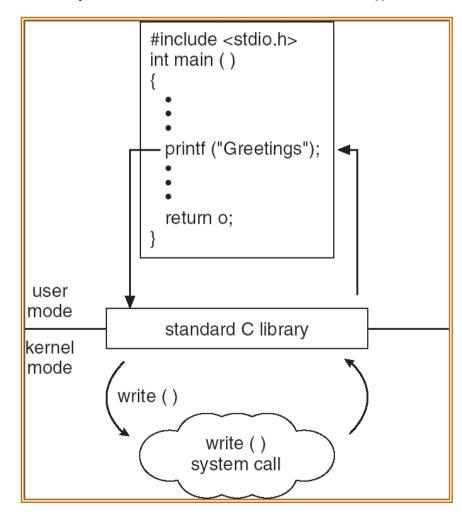
# API – System Call – OS Relationship



# Standard C Library Example

C program invoking printf() library call, which calls write()

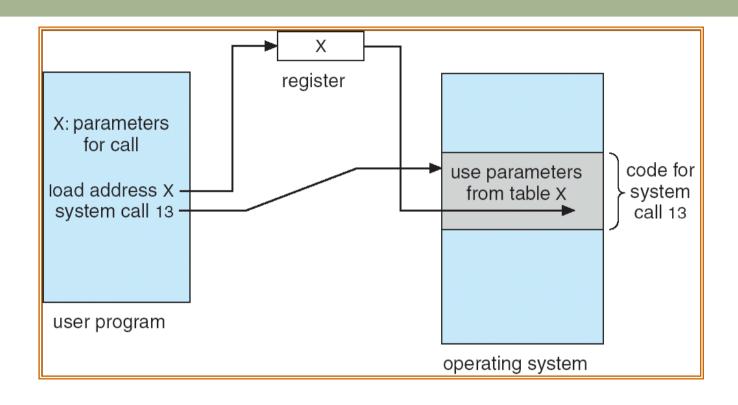
system call



# System Call Parameter Passing

- Three general methods used to pass parameters to the OS
  - Simplest: pass the parameters in *registers* 
    - In some cases, may be more parameters than registers
  - Parameters stored in a block, or table, in memory, and address of block passed as a parameter in a register
    - This approach taken by Linux and Solaris
  - Parameters placed, or *pushed*, onto the *stack* by the program and *popped* off the stack by the operating system
  - Block and stack methods do not limit the number or length of parameters being passed

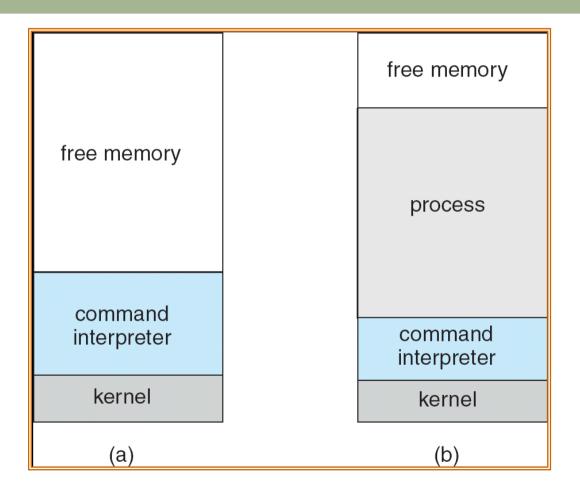
# Parameter Passing via Table



# Types of System Calls

- Process control
- File management
- Device management
- Information maintenance
- Communications

## MS-DOS execution



(a) At system startup (b) running a program

# FreeBSD Running Multiple Programs

process D

free memory

process C

interpreter

process B

kernel

# System Programs

- A convenient environment for program development and execution. They can be divided into:
  - File manipulation
  - Status information
  - File modification
  - Programming language support
  - Program loading and execution
  - Communications
  - Application programs
- Most users' view of the operation system is defined by system programs, not the actual system calls

## System Programs

- Provide a convenient environment for program development and execution
  - Some of them are simply user interfaces to system calls; others are considerably more complex
- File management
  - Create, delete, copy, rename, print, dump, list, and generally manipulate files and directories

# System Programs (cont.)

### Status information

- Some ask the system for info date, time, amount of available memory, disk space, number of users
- Others provide detailed performance, logging, and debugging information
- Typically, these programs format and print the output to the terminal or other output devices
- Some systems implement a registry used to store and retrieve configuration information

# System Programs (cont.)

- File modification
  - Text editors to create and modify files
  - Special commands to search contents of files or perform transformations of the text
- Programming-language support
  - Compilers, assemblers, debuggers and interpreters sometimes provided
- Program loading and execution
  - Absolute loaders, relocatable loaders, linkage editors, and overlay-loaders, debugging systems for higher-level and machine language

# System Programs (cont.)

#### Communications

- Provide the mechanism for creating virtual connections among processes, users, and computer systems
- Allow users to send messages to one another's screens, browse web pages, send electronic-mail messages, log in remotely, transfer files from one machine to another

# Operating System Design and Implementation

- Design and Implementation of OS not "solvable", but some approaches have proven successful
- Internal structure of different Operating Systems can vary widely
- Start by defining goals and specifications
- Affected by choice of hardware, type of system

# Operating System Design and Implementation (Cont.)

- User goals and System goals
  - User goals
    - operating system should be convenient to use, easy to learn, reliable, safe, and fast
  - System goals
    - operating system should be easy to design, implement, and maintain, as well as flexible, reliable, error-free, and efficient
- Important principles to separate
  - **Policy:** What will be done?
  - **Mechanism:** How to do it?

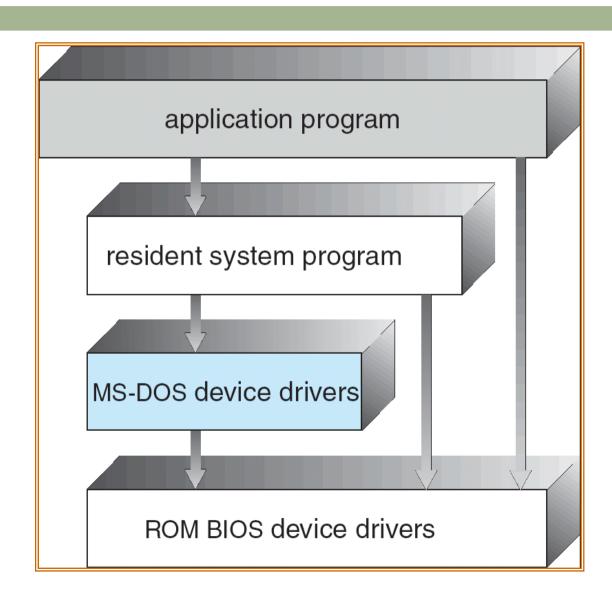
# Operating System Design and Implementation (Cont.)

- Mechanisms determine how to do something, policies decide what will be done
  - The separation of policy from mechanism is a very important principle, it allows maximum flexibility if policy decisions are to be changed later

# Simple Structure

- MS-DOS written to provide the most functionality in the least space
  - Not divided into modules
  - Although MS-DOS has some structure, its interfaces and levels of functionality are not well separated

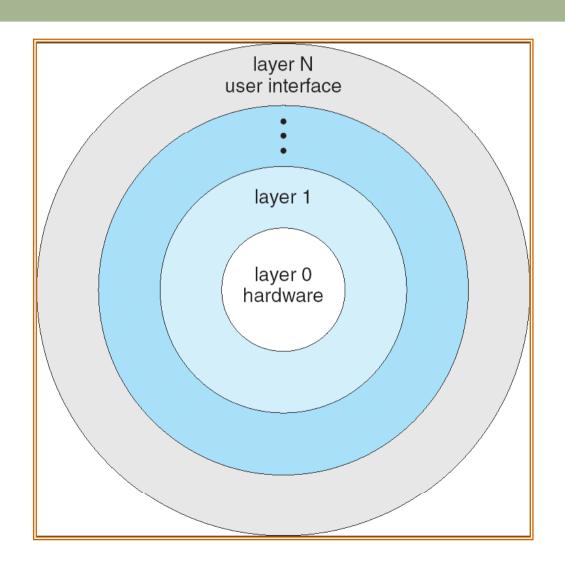
# MS-DOS Layer Structure



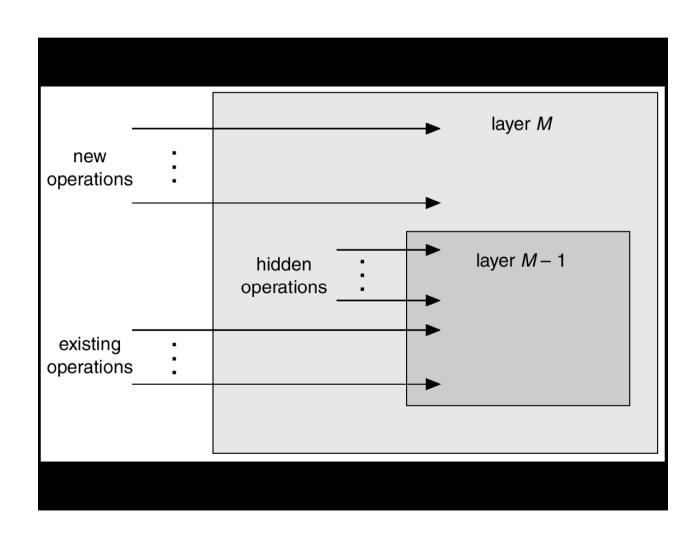
# Layered Approach

- □ The operating system is divided into a number of layers (levels)
  - The bottom layer (layer 0), is the hardware; the highest (layer N) is the user interface.
- With modularity, layers are selected such that each uses functions (operations) and services of only lower-level layers

# Layered Operating System



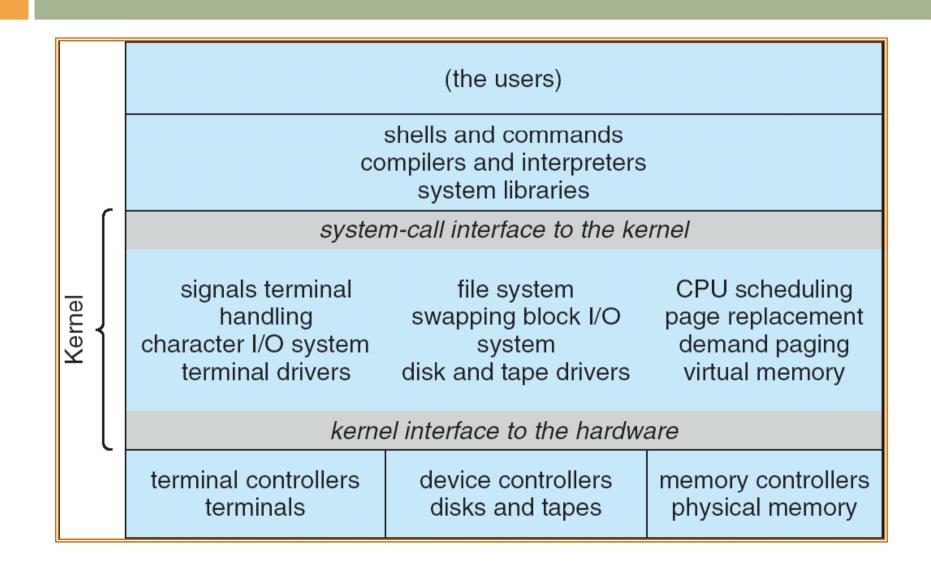
# **Layered Operating System**



## UNIX

- UNIX limited by hardware functionality, the original UNIX operating system had limited structuring.
- The UNIX OS consists of two separable parts
  - Systems programs
  - The kernel
    - Consists of everything below the system-call interface and above the physical hardware
    - Provides the file system, CPU scheduling, memory management, and other operating-system functions; a large number of functions for one level

# **UNIX System Structure**



# Microkernel System Structure

- Moves as much from the kernel into "user" space
- Communication takes place between user modules using message passing

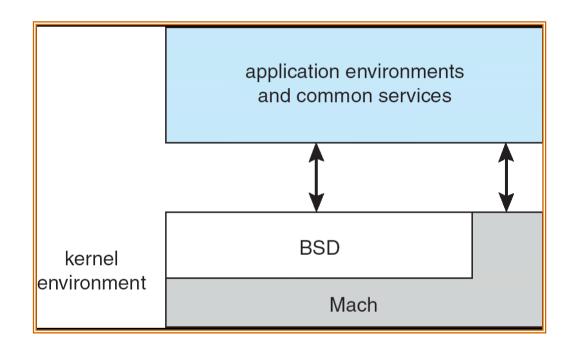
### Benefits:

- Easier to extend a microkernel
- Easier to port the operating system to new architectures
- More reliable (less code is running in kernel mode)
- More secure

### Detriments:

■ Performance overhead of user space to kernel space communication

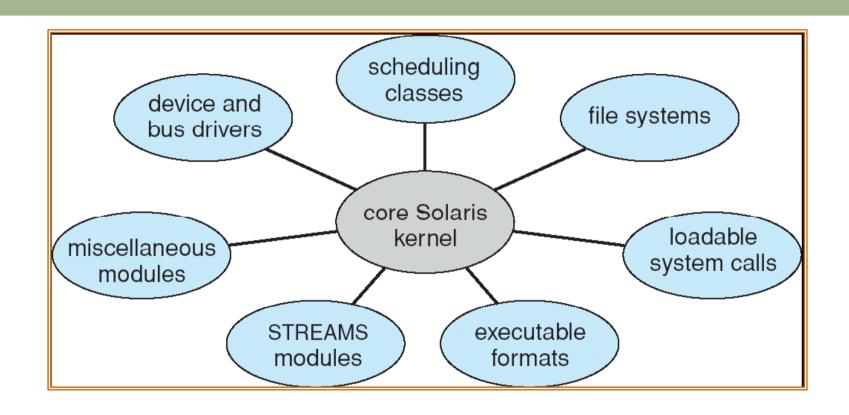
## Mac OS X Structure



### Modules

- Most modern operating systems implement kernel modules
  - Uses object-oriented approach
  - Each core component is separate
  - Each talks to the others over known interfaces
  - Each is loadable as needed within the kernel
- Overall, similar to layers but with more flexible

# Solaris Modular Approach



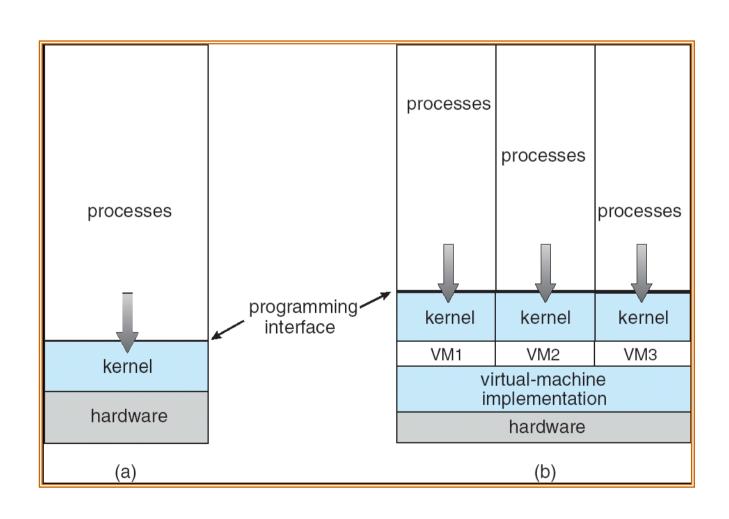
### Virtual Machines

- Takes the layered approach to its logical conclusion.
  - It treats hardware and the operating system kernel as though they were all hardware
- A virtual machine provides an interface identical to the underlying bare hardware
- The operating system creates the illusion of multiple processes.
  - Each executing on its own processor with its own (virtual) memory

# Virtual Machines (Cont.)

- The resources of the physical computer are shared to create the virtual machines
  - CPU scheduling can create the appearance that users have their own processor
  - Spooling and a file system can provide virtual card readers and virtual line printers
  - A normal user time-sharing terminal serves as the virtual machine operator's console
- Difficult to implement due to the effort required to provide an exact duplicate to the underlying machine
  - □ Virtual user mode, virtual monitor mode: in the physical user mode.

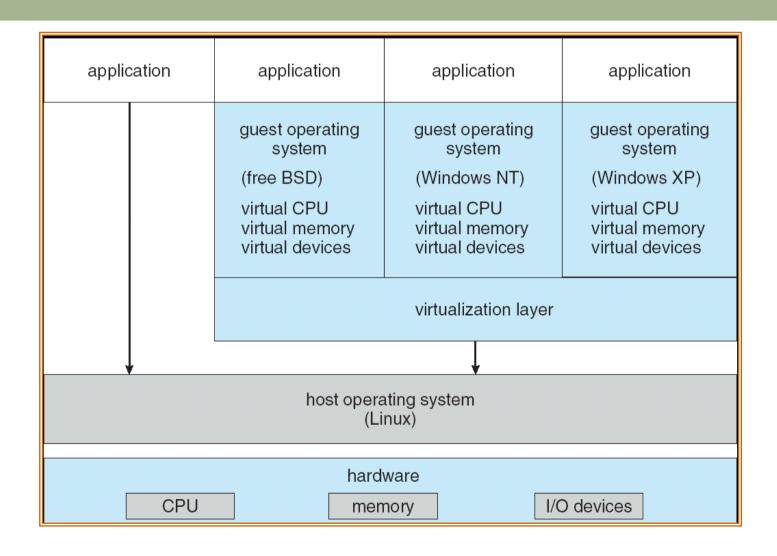
# Virtual Machines (Cont.)



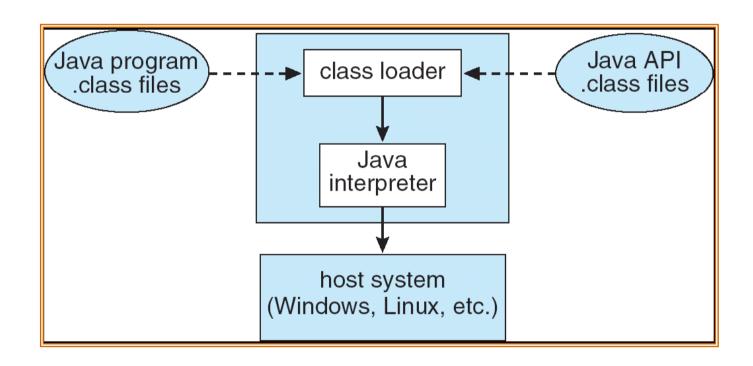
### Benefits of Virtual Machines

- Provides complete protection of system resources
  - Each virtual machine is isolated from all other virtual machines. This isolation, however, permits no direct sharing of resources.
- A perfect vehicle for operating-systems research and development.
  - Instead of on a physical machine; does not disrupt normal system operation.
- A means of solving system compatibility problems

### VMware Architecture



### The Java Virtual Machine



# System Implementation

- Traditionally written in assembly language, operating systems can now be written in higher-level languages. (Unix, WinNT, OS/2)
  - Unix: parts of the scheduler and device drivers are implemented in assembly language.
- Code written in a high-level language:
  - can be written faster.
  - is more compact.
  - is easier to understand and debug.
- An operating system is far easier to port (move to some other hardware) if it is written in a high-level language.

# Operating System Generation

- Operating systems are designed to run on any of a class of machines; the system must be configured for each specific computer site
- Booting starting a computer by loading the kernel
- Bootstrap program code stored in ROM that is able to locate the kernel, load it into memory, and start its execution

# System Boot

- Operating system must be made available to hardware so hardware can start it
  - Small piece of code bootstrap loader, locates the kernel, loads it into memory, and starts it
  - Sometimes two-step process where **boot block** at fixed location loads bootstrap loader
  - When power initialized on system, execution starts at a fixed memory location
    - Firmware used to hold initial boot code
- MBR (Master Boot Record): typically, the 1st sector of the hard disk

# END OF CHAPTER 2